

How to write an efficient SQL

TINGKUN YANG

YUNHE ENMO (BEIJING) TECHNOLOGY CO.,LTD

云和恩墨 成就所托

Introduction

□ 杨廷琨 TINGKUN YANG

- Oracle ACE Director
- ACOUG Vice President
- One of Authors of <Oracle Database Performance Tuning>, <Oracle DBA Notes>, <Oracle DBA Notes 3> & <Oracle Performance Optimization and Diagnosis Case Selection>
- There are more than 2500 technical articles in personal BLOG <http://yangtingkun.net>
- CTO of Enmotech which company is the leader of Database Services industry in CHINA, and include 5 ACE Directors and 4 ACEs and more than 20 Oracle OCM



ORACLE®
ACE Director

ACOUG
All China Oracle User Group
中国 Oracle 用户组

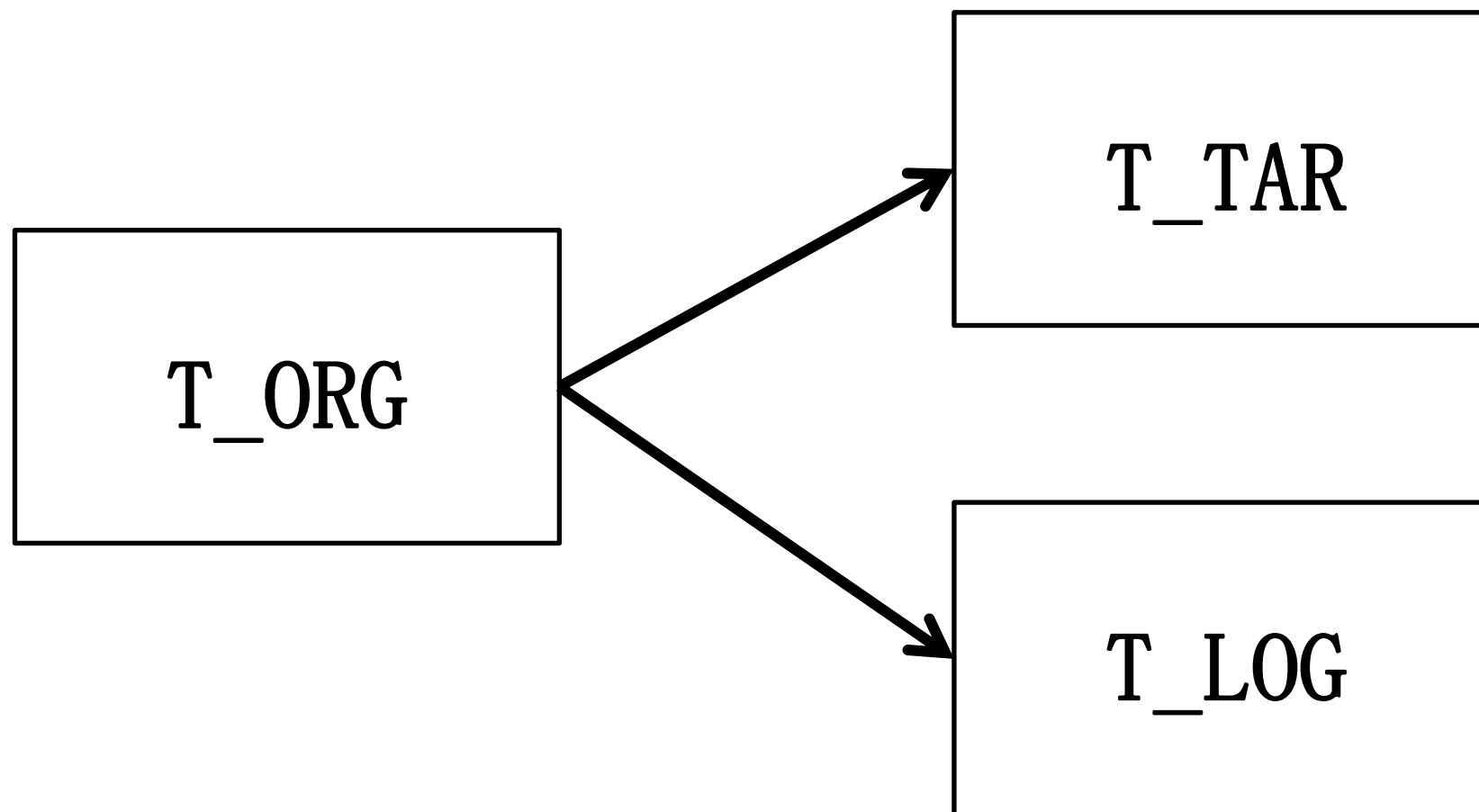
Why do we need efficient SQL

- SQL is currently the second largest programming language, 50% of developers are using SQL
- 80% of database problems are caused by SQL statements
- 80% of the SQL problems come from 20% of SQL statements
- Performance failures caused by a single SQL in high-concurrency systems are very common

- 1 Use new feature reasonably
- 2 Data processing as a whole
- 3 Design SQL execution plan
- 4 Filter unnecessary data from beginning

Use new feature reasonably

Target: We need to read the data from table T_ORG, and insert into table T_TAR, at the same time, we must insert the same data into another table T_LOG.



Use new feature reasonably

Methods:

- Create Trigger

```
CREATE OR REPLACE TRIGGER T_INSERT_TAR
  AFTER INSERT ON T_TAR
  FOR EACH ROW
BEGIN
  INSERT INTO T_LOG (ID, NAME)
  VALUES (:NEW.ID, :NEW.NAME);
END;
/
```

Use new feature reasonably

Methods:

- Create Trigger
 - Too heavy
 - Differences between implementation and demand
 - Increase follow-up maintenance costs
 - Trigger is less efficient

Use new feature reasonably

Methods:

- Create Trigger
- Double Insert

```
INSERT INTO T_TAR SELECT * FROM T_ORG;
```

```
INSERT INTO T_LOG SELECT * FROM T_ORG;
```


Use new feature reasonably

Methods:

- Create Trigger
- Double Insert
 - Consistency
 - Lock
 - Serial transaction
 - Temporary table
 - As of query
 - Atomicity
 - Exception handling

Use new feature reasonably

Methods:

- Create Trigger
- Double Insert
- Open Cursor

SQL> BEGIN

```
2      FOR I IN (SELECT * FROM T_ORG) LOOP
3          INSERT INTO T_TAR VALUES (I.ID, I.NAME);
4          INSERT INTO T_LOG VALUES (I.ID, I.NAME);
5      END LOOP;
6  END;
7  /
```

PL/SQL procedure successfully completed.

Use new feature reasonably

Methods:

- Create Trigger
- Double Insert
- Open Cursor
 - Low efficiency

Use new feature reasonably

Methods:

- Create Trigger
- Double Insert
- Open Cursor
- Bulk Into Variable

Use new feature reasonably

```
SQL> DECLARE
2     TYPE T_ID IS TABLE OF T_ORG.ID%TYPE;
3     TYPE T_NAME IS TABLE OF T_ORG.NAME%TYPE;
4     V_ID T_ID;
5     V_NAME T_NAME;
6 BEGIN
7     SELECT ID, NAME BULK COLLECT INTO V_ID, V_NAME
8     FROM T_ORG;
9     FORALL I IN 1..V_ID.COUNT
10        INSERT INTO T_TAR VALUES(V_ID(I), V_NAME(I));
11     FORALL I IN 1..V_ID.COUNT
12        INSERT INTO T_LOG VALUES(V_ID(I), V_NAME(I));
13 END;
14 /
```

PL/SQL procedure successfully completed.

Use new feature reasonably

Methods:

- Create Trigger
- Double Insert
- Open Cursor
- Bulk Into Variable
 - High complexity
 - Less efficient

Use new feature reasonably

Methods:

- Create Trigger
- Double Insert
- Open Cursor
- Bulk Into Variable
- **Insert All**

```
SQL> INSERT ALL INTO T_TAR (ID, NAME)
2    INTO T_LOG (ID, NAME)
3    SELECT ID, NAME FROM T_ORG;
```

Use new feature reasonably

“New” feature:

- INSERT ALL/ANY Insert multiple tables or insert with condition
- MERGE Combining UPDATE and INSERT statements
- WITH Reduce reading of duplicate result sets
- ANALYTIC FUNCTIONS Row-level calculations avoid self-join
- TOP-N Pagination query
- PIVOT/UNPIVOT Rank conversion query
- MODEL Report array processing
- PATTERN Process complex patterns contained in the data

- 1 Use new feature reasonably
- 2 Data processing as a whole
- 3 Design SQL execution plan
- 4 Filter unnecessary data from beginning

Data processing as a whole

Determine whether the tablespace is automatically expanded:

```
SQL> select distinct tablespace_name, autoextensible
2   from DBA_DATA_FILES
3   where autoextensible = 'YES'
4   union
5   select distinct tablespace_name, autoextensible
6   from DBA_DATA_FILES
7   where autoextensible = 'NO'
8   and tablespace_name not in
9   (select distinct tablespace_name
10   from DBA_DATA_FILES
11   where autoextensible = 'YES');
```

TABLESPACE_NAME	AUT
SYSAUX	YES
SYSTEM	YES
TEST	NO
UNDOTBS1	YES
USERS	YES

Data processing as a whole

```
SQL> select distinct tablespace_name, autoextensible
2   from DBA_DATA_FILES
3   where autoextensible = 'YES'
4   union
5   select distinct tablespace_name, autoextensible
6   from DBA_DATA_FILES
7   where autoextensible = 'NO'
8     and tablespace_name not in
9     (select distinct tablespace_name
10      from DBA_DATA_FILES
11      where autoextensible = 'YES');
```

- benefits: clear thinking
- Disadvantages: Inefficient and redundant
 - Scan view DBA_DATA_FILES three times
 - Include three DISTINCT operation
 - Include UNION data set operation
 - Include NOT IN subquery

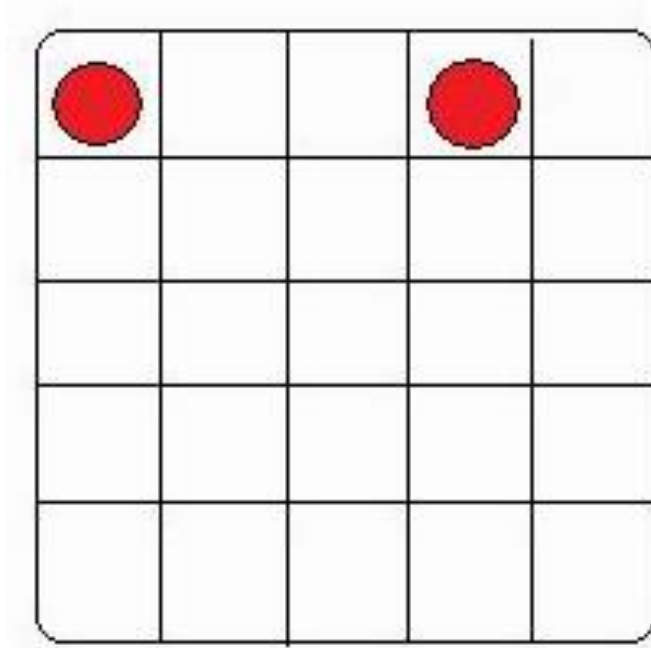
Data processing as a whole

```
SQL> SELECT TABLESPACE_NAME, MAX(AUTOEXTENSIBLE)
2  FROM DBA_DATA_FILES
3  GROUP BY TABLESPACE_NAME;
```

TABLESPACE_NAME	MAX
-----	----
SYSAUX	YES
UNDOTBS1	YES
USERS	YES
TEST	NO
SYSTEM	YES

For strings, 'YES' is bigger than 'NO'.

Data processing as a whole



Question: In a 5x5 chessboard, each row, column, and slash (the slashes not only include diagonal lines) can hold up to two balls. In order to place as many balls as possible, how many ways do you have in total?

Demand: Implementation with one SQL statement.

Output format: From the first row to the fifth row of the checkerboard, each row is sequentially output from the first column to the fifth column.

0 means not to put the ball and 1 means to put. For example :

10010000000000000000000000.

Data processing as a whole

```
SQL> with c as
2   (select rownum - 1 c from dual connect by rownum <= 2),
3   line_seed as
4   (select c1.c || c2.c || c3.c || c4.c || c5.c line,
5    c1.c c1, c2.c c2, c3.c c3, c4.c c4, c5.c c5
6    from c c1, c c2, c c3, c c4, c c5
7    where c1.c + c2.c + c3.c + c4.c + c5.c = 2
8    order by 1 desc)
9   select * from line_seed;
```

LINE	C1	C2	C3	C4	C5
11000	1	1	0	0	0
10100	1	0	1	0	0
10010	1	0	0	1	0
10001	1	0	0	0	1
01100	0	1	1	0	0
01010	0	1	0	1	0
01001	0	1	0	0	1
00110	0	0	1	1	0
00101	0	0	1	0	1
00011	0	0	0	1	1

10 rows selected.

Data processing as a whole

```
with c as
(select rownum - 1 c from dual connect by rownum <= 2),
line_seed as
(select c1.c || c2.c || c3.c || c4.c || c5.c line, c1.c c1, c2.c c2, c3.c c3, c4.c c4, c5.c c5
from c c1, c c2, c c3, c c4, c c5
where c1.c + c2.c + c3.c + c4.c + c5.c = 2
order by 1 desc)
select rownum, line1.line || line2.line || line3.line || line4.line || line5.line result
from line_seed line1, line_seed line2, line_seed line3, line_seed line4, line_seed line5
where line1.c1 + line2.c1 + line3.c1 + line4.c1 + line5.c1 = 2
and line1.c2 + line2.c2 + line3.c2 + line4.c2 + line5.c2 = 2
and line1.c3 + line2.c3 + line3.c3 + line4.c3 + line5.c3 = 2
and line1.c4 + line2.c4 + line3.c4 + line4.c4 + line5.c4 = 2
and line1.c5 + line2.c5 + line3.c5 + line4.c5 + line5.c5 = 2
and line3.c1 + line4.c2 + line5.c3 <= 2
and line2.c1 + line3.c2 + line4.c3 + line5.c4 <= 2
and line1.c1 + line2.c2 + line3.c3 + line4.c4 + line5.c5 <= 2
and line1.c2 + line2.c3 + line3.c4 + line4.c5 <= 2
and line1.c3 + line2.c4 + line3.c5 <= 2
and line1.c3 + line2.c2 + line3.c1 <= 2
and line1.c4 + line2.c3 + line3.c2 + line4.c1 <= 2
and line1.c5 + line2.c4 + line3.c3 + line4.c2 + line5.c1 <= 2
and line2.c5 + line3.c4 + line4.c3 + line5.c2 <= 2
and line3.c5 + line4.c4 + line5.c3 <= 2;
```

Data processing as a whole

```
SQL> with c as
  2  (select rownum - 1 c from dual connect by rownum <= 2),
  3  lines as
  4  (select to_number(c1.c || c2.c || c3.c || c4.c || c5.c) line
  5  from c c1, c c2, c c3, c c4, c c5
  6  where c1.c + c2.c + c3.c + c4.c + c5.c = 2
  7  order by 1 desc)
  8  select ltrim(to_char(line1.line, '09999')) || chr(10)
  9      || ltrim(to_char(line2.line, '09999')) || chr(10)
 10      || ltrim(to_char(line3.line, '09999')) || chr(10)
 11      || ltrim(to_char(line4.line, '09999')) || chr(10)
 12      || ltrim(to_char(line5.line, '00009')) result
 13  from lines line1, lines line2, lines line3, lines line4, lines line5
 14  where line1.line + line2.line + line3.line + line4.line + line5.line = 22222
 15  and ltrim(to_char(line1.line + 10*line2.line + 100*line3.line + 1000*line4.line +
 10000*line5.line), '012') is null
 16  and ltrim(to_char(10000*line1.line + 1000*line2.line + 100*line3.line + 10*line4.line +
  line5.line), '012') is null
 17  and rownum = 1;
```

RESULT

```
-----
11000
10010
00011
01100
00101
```


Data processing as a whole

where $\text{line1.line} + \text{line2.line} + \text{line3.line} + \text{line4.line} + \text{line5.line} = 22222$
 and $\text{ltrim}(\text{to_char}(\text{line1.line} + 10*\text{line2.line} + 100*\text{line3.line} + 1000*\text{line4.line} + 10000*\text{line5.line}), '012')$ is null
 and $\text{ltrim}(\text{to_char}(10000*\text{line1.line} + 1000*\text{line2.line} + 100*\text{line3.line} + 10*\text{line4.line} + \text{line5.line}), '012')$ is null

RESULT

+	1	1	0	0	0
	1	0	0	1	0
	0	0	0	1	1
	0	1	1	0	0
	0	0	1	0	1

	2	2	2	2	2

Data processing as a whole

where $\text{line1.line} + \text{line2.line} + \text{line3.line} + \text{line4.line} + \text{line5.line} = 22222$
 and $\text{ltrim}(\text{to_char}(\text{line1.line} + 10*\text{line2.line} + 100*\text{line3.line} + 1000*\text{line4.line} + 10000*\text{line5.line}), '012')$ is null
 and $\text{ltrim}(\text{to_char}(10000*\text{line1.line} + 1000*\text{line2.line} + 100*\text{line3.line} + 10*\text{line4.line} + \text{line5.line}), '012')$ is null

RESULT

+	1	1	0	0	0				
	1	0	0	1	0				
	0	0	0	0	0				
	0	1	0	0	0				
	0	0	0	0	0				
<hr/>									
	0	0	2	2	2	2	2	0	0

Data processing as a whole

where $\text{line1.line} + \text{line2.line} + \text{line3.line} + \text{line4.line} + \text{line5.line} = 22222$
 and $\text{ltrim}(\text{to_char}(\text{line1.line} + 10*\text{line2.line} + 100*\text{line3.line} + 1000*\text{line4.line} + 10000*\text{line5.line}), '012')$ is null
 and $\text{ltrim}(\text{to_char}(10000*\text{line1.line} + 1000*\text{line2.line} + 100*\text{line3.line} + 10*\text{line4.line} + \text{line5.line}), '012')$ is null

RESULT

+	1	1	0	0	0	0	0	0
		1	0	0	1	0	0	0
			0	0	0	1	1	0
				0	1	1	0	0
					0	0	1	0
<hr/>								
	1	2	0	0	2	2	2	0
								1

Data processing as a whole

- Tom said:
 - You should do it in a single SQL statement if at all possible
 - If you cannot do it in a single SQL Statement, then do it in PL/SQL
 - If you cannot do it in PL/SQL, try a Java Stored Procedure.
- We said:
 - Do not use multiple SQL if you can use single.
 - Do not scan multiple times if you can scan once.
 - Processing the data as a whole not to process row by row.

- 1 Use new feature reasonably
- 2 Data processing as a whole
- 3 Design SQL execution plan
- 4 Filter unnecessary data from beginning

Design SQL execution plan

A customer's SQL statement is slow to execute and consumes a lot of logical reads

```
SQL> EXPLAIN PLAN FOR
 2  SELECT VCM.POL_ID, TC.MOBILE, TC.FIRST_NAME PH_NAME,
 3         PKG_UTILS.F_GET_DESC('T_TITLE', DECODE(TC.GENDER, 'M', '1', 'F', '2'), '211'),
 4         PKG_UTILS.F_GET_DESC('V_PRO_LIFE_3',
 5         (SELECT T.PROD_ID FROM T_CON_PROD T WHERE T.POL_ID = VCM.POL_ID AND T.MASTER_ID IS NULL), '211'),
 6         SUM(TPA.FEE_AMOUNT), VCM.POLICY_CODE, PKG_UTILS.F_GET_DESC('T_LIA_STAT', VCM.LIA_STATE, '211'),
 7         PKG_UTILS.F_GET_DESC('T_SAL_CHAN', VCM.CHANN_TYPE, '211'), VCM.CHANN_TYPE, TC.CUSTOMER_ID, TCO.ORGAN_CODE
 8  FROM V_CON_MAS VCM, T_CON_MAS TCM, T_AGENT TA, T_CUSTOMER TC,
 9         T_POH VPH, T_COM_ORG TCO, T_P_A TPA
10  WHERE VCM.POL_ID = VPH.POL_ID
11  AND VCM.SERVICE_AGENT = TA.AGENT_ID
12  AND VCM.POL_ID = TCM.POL_ID
13  AND VPH.PARTY_ID = TC.CUSTOMER_ID
14  AND VCM.ORGAN_ID = TCO.ORGAN_ID
15  AND VCM.POL_ID = TPA.POL_ID
16  AND VCM.LIABILITY_STATE = 1
17  AND TPA.FEE_TYPE IN (41, 569)
18  AND TPA.FEE_STATUS <> 2
19  AND (EXISTS (
20      SELECT 1 FROM T_PO_CH T, T_CH TC WHERE T.MAS_C_ID = TC.CH_ID AND TC.CH_ID = TPA.CH_ID AND T.SERV_ID = 3 ))
21  AND TPA.DUE_TIME >= (TRUNC(:B1) + 7)
22  AND TPA.DUE_TIME < (TRUNC(:B1) + 8)
23  AND REGEXP_LIKE(TRIM(TC.MOBILE), '^\\d{11}$')
24  AND NOT EXISTS (
25      SELECT 1 FROM T_DATA_EXT SDE, T_DATA TSD
26      WHERE SDE.DATA_ID = TSD.DATA_ID AND SDE.RELAT_VALUE_1 = VCM.POL_ID AND TSD.SMS_ID = 12 AND SDE.RELAT_VALUE_2
= TO_CHAR(:B1, 'yyyy-MM-dd'))
27  GROUP BY VCM.POL_ID, TC.MOBILE, TC.FIRST_NAME, TC.GENDER, VCM.POLICY_CODE, VCM.LIABILITY_STATE, VCM.CHANNEL_TYPE
, TC.CUSTOMER_ID, TCO.ORGAN_CODE;
```

Design SQL execution plan

Id	Operation	Name	Rows	Bytes	Cost (%CPU)
0	SELECT STATEMENT		356	53400	901K (1)
* 1	INDEX RANGE SCAN	IDX_CON_PROD__PRD_ID	1	12	1 (0)
2	HASH GROUP BY		356	53400	901K (1)
* 3	FILTER				
* 4	FILTER				
5	NESTED LOOPS		356	53400	84579 (2)
* 6	HASH JOIN		356	51264	84578 (2)
7	TABLE ACCESS FULL	T_C_O	143	2145	4 (0)
8	NESTED LOOPS SEMI		356	45924	84573 (2)
9	NESTED LOOPS		370	44030	84481 (2)
10	NESTED LOOPS		6067	539K	82962 (2)
11	NESTED LOOPS		6036	465K	81452 (2)
* 12	HASH JOIN		6037	430K	81451 (2)
* 13	TABLE ACCESS BY INDEX ROWID	T_P_A	4812	145K	596 (1)
* 14	INDEX RANGE SCAN	IDX_P_A__DUE	14477		10 (0)
15	VIEW	V_CON_MAS	877K	35M	80844 (2)
16	UNION-ALL				
* 17	TABLE ACCESS FULL	T_CON_MAS_LOG	877K	36M	80844 (2)
* 18	FILTER				
* 19	TABLE ACCESS FULL	T_CON_MAS	155K	6357K	21715 (1)
* 20	INDEX UNIQUE SCAN	PK_T_AGENT	1	6	1 (0)
21	TABLE ACCESS BY INDEX ROWID	T_POH	1	12	1 (0)
* 22	INDEX RANGE SCAN	IDX_POH_PO_ID	1		1 (0)
* 23	TABLE ACCESS BY INDEX ROWID	T_CUSTOMER	1	28	1 (0)
* 24	INDEX UNIQUE SCAN	PK_T_CUSTOMER	1		1 (0)
* 25	TABLE ACCESS BY INDEX ROWID	T_PO_CH	244K	2389K	1 (0)
* 26	INDEX RANGE SCAN	IDX_PO_CH__MAS_CH_ID	1		1 (0)
* 27	INDEX UNIQUE SCAN	PK_T_CON_MAS	1	6	1 (0)
28	NESTED LOOPS				
29	NESTED LOOPS		1	34	4591 (1)
30	TABLE ACCESS BY INDEX ROWID	T_DATA_EXT	1	25	4590 (1)
* 31	INDEX SKIP SCAN	IDX_DATA_EXT__RELAT_	1		4589 (1)
* 32	INDEX UNIQUE SCAN	PK_T_DATA	1		1 (0)
* 33	TABLE ACCESS BY INDEX ROWID	T_DATA	1	9	1 (0)

Design SQL execution plan

```
SQL> create table t_objects as select * from dba_objects;
```

Table created.

```
SQL> create index ind_obj_id on t_objects(object_id);
```

Index created.

```
SQL> create table t_tab as select * from dba_tables;
```

Table created.

```
SQL> create table t_ind as select * from dba_indexes;
```

Table created.

```
SQL> create index ind_tab_name on t_tab(table_name);
```

Index created.

```
SQL> create index ind_ind_name on t_ind(index_name);
```

Index created.

```
SQL> create view v_seg as
```

```
2 select owner, table_name, tablespace_name, blocks from t_tab where temporary = 'N'
```

```
3 union all
```

```
4 select owner, index_name, tablespace_name, num_rows from t_ind where status = 'N/A';
```

View created.

Design SQL execution plan

```
SQL> select obj.owner, obj.object_name, created, v.blocks  
2  from t_objects obj, v_seg v  
3  where obj.object_id = 12345  
4  and obj.object_name = v.table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	194	43
* 1	HASH JOIN		1	194	43
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	115	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW	V_SEG	3340	257K	40
5	UNION-ALL				
* 6	TABLE ACCESS FULL	T_TAB	1184	95904	15
* 7	TABLE ACCESS FULL	T_IND	2156	178K	25

Design SQL execution plan

```
SQL> select /*+ index(t_tab ind_tab_name) */ obj.owner, obj.object_name, created, v.blocks  
2  from t_objects obj, v_seg v  
3  where obj.object_id = 12345  
4  and obj.object_name = v.table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	194	43
* 1	HASH JOIN		1	194	43
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	115	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW	V_SEG	3340	257K	40
5	UNION-ALL				
* 6	TABLE ACCESS FULL	T_TAB	1184	95904	15
* 7	TABLE ACCESS FULL	T_IND	2156	178K	25

Design SQL execution plan

```
SQL> select /*+ index(v.t_tab ind_tab_name) */ obj.owner, obj.object_name, created, v.blocks  
2  from t_objects obj, v_seg v  
3  where obj.object_id = 12345  
4  and obj.object_name = v.table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	194	1210
* 1	HASH JOIN		1	194	1210
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	115	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW	V_SEG	3340	257K	1207
5	UNION-ALL				
* 6	TABLE ACCESS BY INDEX ROWID	T_TAB	1184	95904	1182
7	INDEX FULL SCAN	IND_TAB_NAME	2367		11
* 8	TABLE ACCESS FULL	T_IND	2156	178K	25

Design SQL execution plan

```
SQL> select obj.owner, obj.object_name, created, v.blocks
2  from t_objects obj,
3  (select owner, table_name, tablespace_name, blocks from t_tab where temporary = 'N'
4  union all
5  select owner, index_name, tablespace_name, num_rows from t_ind where status = 'N/A') v
6  where obj.object_id = 12345
7  and obj.object_name = v.table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	194	43
* 1	HASH JOIN		1	194	43
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	115	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW		3340	257K	40
5	UNION-ALL				
* 6	TABLE ACCESS FULL	T_TAB	1184	95904	15
* 7	TABLE ACCESS FULL	T_IND	2156	178K	25

Design SQL execution plan

```
SQL> select /*+ index(v.t_tab ind_tab_name) */ obj.owner, obj.object_name, created, v.blocks
2  from t_objects obj,
3  (select owner, table_name, tablespace_name, blocks from t_tab where temporary = 'N'
4  union all
5  select owner, index_name, tablespace_name, num_rows from t_ind where status = 'N/A') v
6  where obj.object_id = 12345
7  and obj.object_name = v.table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	194	1210
* 1	HASH JOIN		1	194	1210
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	115	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW		3340	257K	1207
5	UNION-ALL				
* 6	TABLE ACCESS BY INDEX ROWID	T_TAB	1184	95904	1182
7	INDEX FULL SCAN	IND_TAB_NAME	2367		11
* 8	TABLE ACCESS FULL	T_IND	2156	178K	25

Design SQL execution plan

```
SQL> select /*+ index(t_tab ind_tab_name) */ obj.owner, obj.object_name, created, blocks
2  from t_objects obj,
3  (select owner, table_name, tablespace_name, blocks from t_tab where temporary = 'N'
4  union all
5  select owner, index_name, tablespace_name, num_rows from t_ind where status = 'N/A')
6  where obj.object_id = 12345
7  and obj.object_name = table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	194	43
* 1	HASH JOIN		1	194	43
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	115	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW		3340	257K	40
5	UNION-ALL				
* 6	TABLE ACCESS FULL	T_TAB	1184	95904	15
* 7	TABLE ACCESS FULL	T_IND	2156	178K	25

Design SQL execution plan

```
SQL> SELECT ID, OBJECT_ALIAS, DEPTH
2 FROM V$SQL_PLAN
3 WHERE SQL_ID IN
4 (SELECT SQL_ID FROM V$SQL
5 WHERE SQL_TEXT LIKE 'select /*+ index(t_tab ind_tab_name) */%')
6 ORDER BY SQL_ID, ID;
```

ID	OBJECT_ALIAS	DEPTH
0		0
1		1
2	OBJ@SEL\$1	2
3	OBJ@SEL\$1	3
4	from\$_subquery\$_002@SEL\$1	2
5		3
6	T_TAB@SEL\$2	4
7	T_IND@SEL\$3	4

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	194	43
* 1	HASH JOIN		1	194	43
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	115	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW		3340	257K	40
5	UNION-ALL				
* 6	TABLE ACCESS FULL	T_TAB	1184	95904	15
* 7	TABLE ACCESS FULL	T_IND	2156	178K	25

Design SQL execution plan

```
SQL> select /*+ index("from$_subquery$_002".t_tab ind_tab_name) */ obj.owner, obj.object_name,  
created, blocks  
2 from t_objects obj,  
3 (select owner, table_name, tablespace_name, blocks from t_tab where temporary = 'N'  
4 union all  
5 select owner, index_name, tablespace_name, num_rows from t_ind where status = 'N/A')  
6 where obj.object_id = 12345  
7 and obj.object_name = table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	194	1210
* 1	HASH JOIN		1	194	1210
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	115	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW		3340	257K	1207
5	UNION-ALL				
* 6	TABLE ACCESS BY INDEX ROWID	T_TAB	1184	95904	1182
7	INDEX FULL SCAN	IND_TAB_NAME	2367		11
* 8	TABLE ACCESS FULL	T_IND	2156	178K	25

Design SQL execution plan

```
SQL> select /*+ index(@v_1 t_tab ind_tab_name) */ obj.owner, obj.object_name, created, blocks
  2  from t_objects obj,
  3  (select /*+ qb_name(v_1) */ owner, table_name, tablespace_name, blocks from t_tab where
temporary = 'N'
  4  union all
  5  select owner, index_name, tablespace_name, num_rows from t_ind where status = 'N/A')
  6  where obj.object_id = 12345
  7  and obj.object_name = table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost
0	SELECT STATEMENT		1	123	1210
* 1	HASH JOIN		1	123	1210
2	TABLE ACCESS BY INDEX ROWID	T_OBJECTS	1	44	2
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1
4	VIEW		3340	257K	1207
5	UNION-ALL				
* 6	TABLE ACCESS BY INDEX ROWID	T_TAB	1184	28416	1182
7	INDEX FULL SCAN	IND_TAB_NAME	2367		11
* 8	TABLE ACCESS FULL	T_IND	2156	66836	25

Design SQL execution plan

```
SQL> select /*+ push_pred(v) */ obj.owner, obj.object_name, created, v.blocks
2  from t_objects obj, v_seg v
3  where obj.object_id = 12345
4  and obj.object_name = v.table_name;
```

no rows selected

Id	Operation	Name	Rows	Bytes	Cost (%CPU)
0	SELECT STATEMENT		1	78	6 (0)
1	NESTED LOOPS		1	78	6 (0)
2	TABLE ACCESS BY INDEX ROWID BATCHED	T_OBJECTS	1	44	2 (0)
* 3	INDEX RANGE SCAN	IND_OBJ_ID	1		1 (0)
4	VIEW	V_SEG	1	34	4 (0)
5	UNION ALL PUSHED PREDICATE				
* 6	TABLE ACCESS BY INDEX ROWID BATCHED	T_TAB	1	24	2 (0)
* 7	INDEX RANGE SCAN	IND_TAB_NAME	1		1 (0)
* 8	TABLE ACCESS BY INDEX ROWID BATCHED	T_IND	1	31	2 (0)
* 9	INDEX RANGE SCAN	IND_IND_NAME	1		1 (0)

Design SQL execution plan

```
SQL> alter session set "_optimizer_push_pred_cost_based"=false;
```

Session altered.

Id	Operation	Name	Rows	Bytes	Cost (%CPU)
0	SELECT STATEMENT		286	42900	659K (1)
* 1	INDEX RANGE SCAN	IDX_CON_PROD__PRD_ID	1	12	1 (0)
2	HASH GROUP BY		286	42900	659K (1)
.					
.					
.					
* 12	TABLE ACCESS BY INDEX ROWID	T_P_A	4812	145K	596 (1)
* 13	INDEX RANGE SCAN	IDX_P_A__DUE	14477		10 (0)
* 14	VIEW	V_CON_MAS	1	42	1 (0)
15	UNION-ALL PARTITION				
* 16	TABLE ACCESS BY INDEX ROWID	T_CON_MAS_LOG	1	44	1 (0)
* 17	INDEX RANGE SCAN	IDX_CON_MAS_L__LAST_CM	1		1 (0)
* 18	FILTER				
* 19	TABLE ACCESS BY INDEX ROWID	T_CON_MAS	1	42	1 (0)
* 20	INDEX UNIQUE SCAN	PK_T_CON_MAS	1		1 (0)
21	NESTED LOOPS				
22	NESTED LOOPS		1	34	4591 (1)
23	TABLE ACCESS BY INDEX ROWID	T_DATA_EXT	1	25	4590 (1)
.					
.					
.					
* 34	INDEX UNIQUE SCAN	PK_T_CON_MAS	1	6	1 (0)
* 35	INDEX UNIQUE SCAN	PK_T_COM_ORG	1		1 (0)
36	TABLE ACCESS BY INDEX ROWID	T_COM_ORG	1	15	1 (0)

Design SQL execution plan

The importance of use the HINT:

- Control SQL execution
- Understand the difference in the efficiency of the execution plan
- A powerful tool for analyzing SQL problems
- Bypass bugs or performance issues
- Force SQL to execute according to design ideas

- 1 Use new feature reasonably
- 2 Data processing as a whole
- 3 Design SQL execution plan
- 4 Filter unnecessary data from beginning

Filter unnecessary data from beginning

How to find all prime numbers within 100 through one SQL:

```
SQL> WITH T
  2 AS
  3 (SELECT ROWNUM RN FROM DUAL CONNECT BY LEVEL < 100)
  4 SELECT RN FROM T
  5 WHERE RN > 1
  6 MINUS
  7 SELECT A.RN * B.RN FROM T A, T B
  8 WHERE A.RN <= B.RN
  9 AND A.RN > 1
 10 AND B.RN > 1;
```

```
      RN
-----
      2
      3
      5
.
.
.
     97
```

25 rows selected.

Elapsed: 00:00:00.09

Filter unnecessary data from beginning

```
SQL> WITH T
  2 AS
  3 (SELECT ROWNUM RN FROM DUAL CONNECT BY LEVEL < 10000)
  4 SELECT RN FROM T
  5 WHERE RN > 1
  6 MINUS
  7 SELECT A.RN * B.RN FROM T A, T B
  8 WHERE A.RN <= B.RN
  9 AND A.RN > 1
 10 AND B.RN > 1;
```

1229 rows selected.

Elapsed: 00:01:12.50

Statistics

543	recursive calls
148	db block gets
170002	consistent gets
69148	physical reads
884	redo size
21808	bytes sent via SQL*Net to client
1415	bytes received via SQL*Net from client
83	SQL*Net roundtrips to/from client
2	sorts (memory)
1	sorts (disk)
1229	rows processed

Filter unnecessary data from beginning

```
SQL> WITH T
  2 AS
  3 (SELECT ROWNUM RN FROM DUAL CONNECT BY LEVEL < 10000)
  4 SELECT RN FROM T
  5 WHERE RN > 1
  6 MINUS
  7 SELECT A.RN * B.RN FROM T A, T B
  8 WHERE A.RN <= B.RN
  9 AND A.RN > 1
 10 AND A.RN <= 100
 11 AND B.RN > 1
 12 AND B.RN <= 5000;
```

1229 rows selected.

Elapsed: 00:00:00.52

Statistics

```
      9 recursive calls
     24 db block gets
    1727 consistent gets
     17 physical reads
     604 redo size
   21808 bytes sent via SQL*Net to client
   1415 bytes received via SQL*Net from client
      83 SQL*Net roundtrips to/from client
       3 sorts (memory)
       0 sorts (disk)
    1229 rows processed
```


Filter unnecessary data from beginning

```
SQL> WITH T
  2 AS
  3 (SELECT ROWNUM * 2 + 1 RN FROM DUAL CONNECT BY LEVEL < 4999)
  4 SELECT 2 FROM DUAL
  5 UNION ALL
  6 (
  7  SELECT RN FROM T
  8  MINUS
  9  SELECT A.RN * B.RN FROM T A, T B
 10  WHERE A.RN <= B.RN
 11  AND A.RN <= 100
 12  AND B.RN <= 5000
 13 );
```

1229 rows selected.

Elapsed: 00:00:00.13

Statistics

```
      9 recursive calls
     16 db block gets
    469 consistent gets
      9 physical reads
     604 redo size
  21807 bytes sent via SQL*Net to client
   1415 bytes received via SQL*Net from client
      83 SQL*Net roundtrips to/from client
        3 sorts (memory)
        0 sorts (disk)
    1229 rows processed
```

Filter unnecessary data from beginning

```
SQL> WITH T
  2 AS
  3 (SELECT ROWNUM * 2 + 1 RN FROM DUAL CONNECT BY LEVEL < 4999)
  4 SELECT 2 FROM DUAL
  5 UNION ALL
  6 (
  7  SELECT RN FROM T
  8  MINUS
  9  SELECT A.RN * B.RN FROM T A, T B
 10  WHERE A.RN <= B.RN
 11  AND A.RN <= 99
 12  AND B.RN <= 3333
 13  AND A.RN * B.RN < 10000);
```

1229 rows selected.

Elapsed: 00:00:00.05

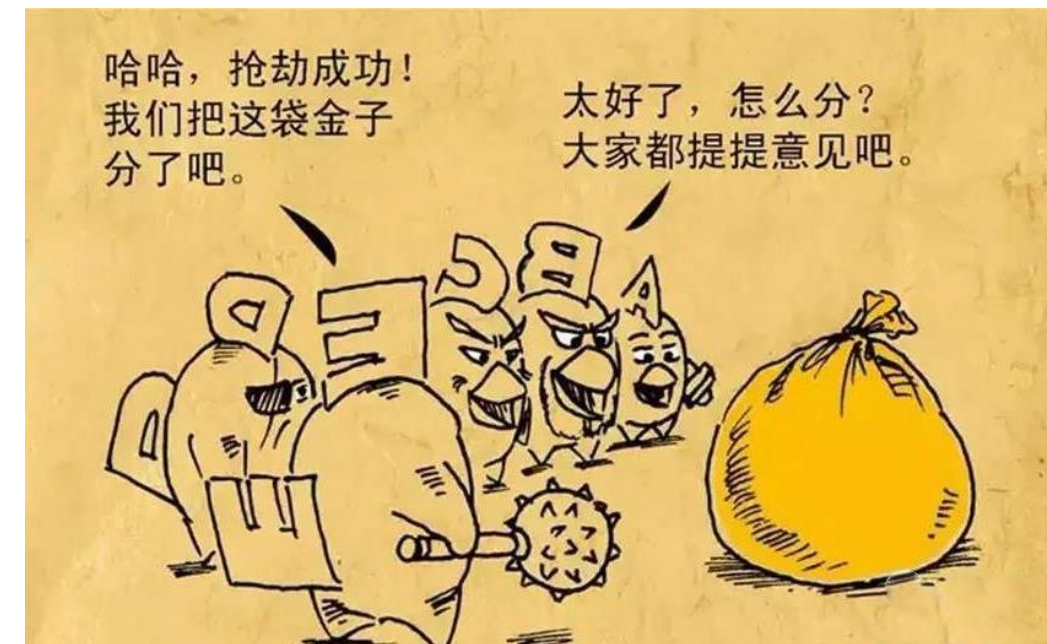
Statistics

```
      2 recursive calls
     19 db block gets
    461 consistent gets
      9 physical reads
     940 redo size
  21807 bytes sent via SQL*Net to client
   1415 bytes received via SQL*Net from client
      83 SQL*Net roundtrips to/from client
       3 sorts (memory)
       0 sorts (disk)
   1229 rows processed
```

Filter unnecessary data from beginning

The issue of how pirate is dividing the gold :

There were five pirates who looted 100 gold coins and wanted to divide the gold coins. The first step they will take randomly from the five notes written from 1 to 5, the number they get will determine their order.



The luckier get 1 can propose a distribution plan. If his plan is approved by more than half of the pirates, he will divide the gold according to his plan, otherwise the first person will be killed. Then the one who get 2 can propose a distribution plan and continue the above rules.

So the question is: If you get 1, what kind of distribution plan would you propose?

(We assume that every pirate is smart enough and pursues maximization of benefits.)

Filter unnecessary data from beginning

Distribution Principle for One: Exclusive

```
SQL> WITH A AS  
2  (SELECT 100 - ROWNUM + 1 N FROM DUAL CONNECT BY ROWNUM <= 101),  
3  MAX_ONE AS  
4  (SELECT MAX(N) MAX1 FROM A)  
5  SELECT * FROM MAX_ONE;
```

MAX1

100

Filter unnecessary data from beginning

Distribution Principle for Two: Sadness

```
SQL> WITH A AS
  2  (SELECT 100 - ROWNUM + 1 N FROM DUAL CONNECT BY ROWNUM <= 101),
  3  MAX_ONE AS
  4  (SELECT MAX(N) MAX1 FROM A),
  5  MAX_TWO AS
  6  (SELECT /*+ LEADING (P2, P1) USE_NL(P1) */ P2.N MAX2, P1.N MAX1
  7  FROM A P1, A P2
  8  WHERE P1.N + P2.N = 100
  9  AND P1.N >= (SELECT MAX1 FROM MAX_ONE)
 10  AND ROWNUM = 1)
 11  SELECT * FROM MAX_TWO;
```

MAX2	MAX1
0	100

Filter unnecessary data from beginning

Distribution Principle for Three: Lure

```
SQL> WITH A AS
  2  (SELECT 100 - ROWNUM + 1 N FROM DUAL CONNECT BY ROWNUM <= 101),
  3  MAX_ONE AS
  4  (SELECT MAX(N) MAX1 FROM A),
  5  MAX_TWO AS
  6  (SELECT /*+ LEADING (P2, P1) USE_NL(P1) */ P2.N MAX2, P1.N MAX1
  7  FROM A P1, A P2
  8  WHERE P1.N + P2.N = 100
  9  AND P1.N >= (SELECT MAX1 FROM MAX_ONE)
10  AND ROWNUM = 1),
11  MAX_THREE AS
12  (SELECT /*+ LEADING(P3, P2, P1) USE_NL(P1) */ P3.N MAX3, P2.N MAX2, P1.N MAX1
13  FROM A P1, A P2, A P3, MAX_TWO
14  WHERE P1.N + P2.N + P3.N = 100
15  AND SIGN(P2.N - MAX2) + SIGN(P1.N - MAX1) >= 0
16  AND ROWNUM = 1)
17  SELECT * FROM MAX_THREE;
```

MAX3	MAX2	MAX1
99	1	0

Filter unnecessary data from beginning

Distribution Principle for Four: Exclude

```
SQL> WITH A AS
  2  (SELECT 100 - ROWNUM + 1 N FROM DUAL CONNECT BY ROWNUM <= 101),
  3  MAX_ONE AS
  4  (SELECT MAX(N) MAX1 FROM A),
  5  MAX_TWO AS
  6  (SELECT /*+ LEADING (P2, P1) USE_NL(P1) */ P2.N MAX2, P1.N MAX1
  7  FROM A P1, A P2
  8  WHERE P1.N + P2.N = 100
  9  AND P1.N >= (SELECT MAX1 FROM MAX_ONE)
10  AND ROWNUM = 1),
11  MAX_THREE AS
12  (SELECT /*+ LEADING(P3, P2, P1) USE_NL(P1) */ P3.N MAX3, P2.N MAX2, P1.N MAX1
13  FROM A P1, A P2, A P3, MAX_TWO
14  WHERE P1.N + P2.N + P3.N = 100
15  AND SIGN(P2.N - MAX2) + SIGN(P1.N - MAX1) >= 0
16  AND ROWNUM = 1),
17  MAX_FOUR AS
18  (SELECT /*+ LEADING(P4, P3, P2, P1) USE_NL(P3) USE_NL(P2) USE_NL(P1) */ P4.N MAX4, P3.N MAX3, P2.N
MAX2, P1.N MAX1
19  FROM A P1, A P2, A P3, A P4, MAX_THREE
20  WHERE P1.N + P2.N + P3.N + P4.N = 100
21  AND SIGN(P3.N - MAX3) + SIGN(P2.N - MAX2) + SIGN(P1.N - MAX1) > 0
22  AND ROWNUM = 1)
23  SELECT * FROM MAX_FOUR;
```

MAX4	MAX3	MAX2	MAX1
97	0	2	1

Filter unnecessary data from beginning

```
SQL> WITH A AS
  2  (SELECT 100 - ROWNUM + 1 N FROM DUAL CONNECT BY ROWNUM <= 101),
  3  MAX_ONE AS
  4  (SELECT MAX(N) MAX1 FROM A),
  5  MAX_TWO AS
  6  (SELECT /*+ LEADING (P2, P1) USE_NL(P1) */ P2.N MAX2, P1.N MAX1
  7  FROM A P1, A P2
  8  WHERE P1.N + P2.N = 100
  9  AND P1.N >= (SELECT MAX1 FROM MAX_ONE)
10  AND ROWNUM = 1),
11  MAX_THREE AS
12  (SELECT /*+ LEADING(P3, P2, P1) USE_NL(P1) */ P3.N MAX3, P2.N MAX2, P1.N MAX1
13  FROM A P1, A P2, A P3, MAX_TWO
14  WHERE P1.N + P2.N + P3.N = 100
15  AND SIGN(P2.N - MAX2) + SIGN(P1.N - MAX1) >= 0
16  AND ROWNUM = 1),
17  MAX_FOUR AS
18  (SELECT /*+ LEADING(P4, P3, P2, P1) USE_NL(P3) USE_NL(P2) USE_NL(P1) */ P4.N MAX4, P3.N MAX3, P2.N MAX2, P1.N MAX1
19  FROM A P1, A P2, A P3, A P4, MAX_THREE
20  WHERE P1.N + P2.N + P3.N + P4.N = 100
21  AND SIGN(P3.N - MAX3) + SIGN(P2.N - MAX2) + SIGN(P1.N - MAX1) > 0
22  AND ROWNUM = 1),
23  FIVE AS
24  (SELECT /*+ LEADING(P5, P4, P3, P2, P1) USE_NL(P4) USE_NL(P3) USE_NL(P2) USE_NL(P1) */ P5.N N5, P4.N N4, P3.N N3,
P2.N N2, P1.N N1
25  FROM A P1, A P2, A P3, A P4, A P5, MAX_FOUR
26  WHERE P1.N + P2.N + P3.N + P4.N + P5.N = 100
27  AND SIGN(P4.N - MAX4) + SIGN(P3.N - MAX3) + SIGN(P2.N - MAX2) + SIGN(P1.N - MAX1) >= 0
28  AND ROWNUM = 1)
29  SELECT * FROM FIVE;
```

N5	N4	N3	N2	N1
97	0	1	0	2

Filter unnecessary data from beginning

```

SQL> WITH A AS
2  (SELECT 100 - ROWNUM + 1 N, ROWNUM - 1 NA FROM DUAL CONNECT BY ROWNUM <= 101),
3  MAX_ONE AS
4  (SELECT MAX(N) MAX1 FROM A),
5  MAX_TWO AS
6  (SELECT /*+ LEADING(MAX_ONE, P2, P1) USE_NL(P2) USE_NL(P1) */ DECODE(P1.NA, MAX1, P2.N - 1, P2.N) MAX2, P1.NA MAX1
7  FROM A P1, A P2, MAX_ONE
8  WHERE P1.NA + P2.N = 100
9  AND P1.NA >= MAX1
10 AND ROWNUM = 1),
11 MAX_THREE AS
12 (SELECT /*+ LEADING(MAX_TWO, P3, P2, P1) USE_NL(P3) USE_NL(P2) USE_NL(P1) */ P3.N MAX3, P2.NA MAX2, P1.NA MAX1
13 FROM A P1, A P2, A P3, MAX_TWO
14 WHERE P1.NA + P2.NA + P3.N = 100
15 AND P3.N + P2.NA <= 100
16 AND CASE WHEN P2.NA > MAX2 THEN 1 ELSE -1 END + CASE WHEN P1.NA > MAX1 THEN 1 ELSE -1 END >= 0
17 AND ROWNUM = 1),
18 MAX_FOUR AS
19 (SELECT /*+ LEADING(MAX_THREE, P4, P3, P2, P1) USE_NL(P4) USE_NL(P3) USE_NL(P2) USE_NL(P1) */ P4.N MAX4, P3.NA MAX3, P2.NA MAX2, P1.NA MAX1
20 FROM A P1, A P2, A P3, A P4, MAX_THREE
21 WHERE P1.NA + P2.NA + P3.NA + P4.N = 100
22 AND P4.N + P3.NA <= 100
23 AND P4.N + P3.NA + P2.NA <= 100
24 AND CASE WHEN P3.NA > MAX3 THEN 1 ELSE -1 END + CASE WHEN P2.NA > MAX2 THEN 1 ELSE -1 END + CASE WHEN P1.NA > MAX1 THEN 1 ELSE -1 END > 0
25 AND ROWNUM = 1),
26 MAX_FIVE AS
27 (SELECT /*+ LEADING(MAX_FOUR, P5, P4, P3, P2, P1) USE_NL(P5) USE_NL(P4) USE_NL(P3) USE_NL(P2) USE_NL(P1) */ P5.N N5, P4.NA N4, P3.NA N3, P2.NA N2, P1.NA N1
28 FROM A P1, A P2, A P3, A P4, A P5, MAX_FOUR
29 WHERE P1.NA + P2.NA + P3.NA + P4.NA + P5.N = 100
30 AND P5.N + P4.NA <= 100
31 AND P5.N + P4.NA + P3.NA <= 100
32 AND P5.N + P4.NA + P3.NA + P2.NA <= 100
33 AND CASE WHEN P4.NA > MAX4 THEN 1 ELSE -1 END + CASE WHEN P3.NA > MAX3 THEN 1 ELSE -1 END + CASE WHEN P2.NA > MAX2 THEN 1 ELSE -1 END + CASE WHEN P1.NA > MAX1 THEN 1 ELSE -1 END >= 0
34 AND ROWNUM = 1),
35 FIVE AS
36 (SELECT /*+ LEADING(MAX_FOUR, MAX_FIVE, P5, P4, P3, P2, P1) USE_NL(MAX_FIVE) USE_NL(P5) USE_NL(P4) USE_NL(P3) USE_NL(P2) USE_NL(P1) */ P5.N N5, P4.NA N4, P3.NA N3, P2.NA N2, P1.NA N1
37 FROM A P1, A P2, A P3, A P4, A P5, MAX_FOUR, MAX_FIVE
38 WHERE P1.NA + P2.NA + P3.NA + P4.NA + P5.N = 100
39 AND MAX_FIVE.N5 = P5.N
40 AND P5.N + P4.NA <= 100
41 AND P5.N + P4.NA + P3.NA <= 100
42 AND P5.N + P4.NA + P3.NA + P2.NA <= 100
43 AND CASE WHEN P4.NA > MAX_FOUR.MAX4 THEN 1 ELSE -1 END
44   + CASE WHEN P3.NA > MAX_FOUR.MAX3 THEN 1 ELSE -1 END
45   + CASE WHEN P2.NA > MAX_FOUR.MAX2 THEN 1 ELSE -1 END
46   + CASE WHEN P1.NA > MAX_FOUR.MAX1 THEN 1 ELSE -1 END >= 0)
47 SELECT * FROM FIVE;

```

N5	N4	N3	N2	N1
97	0	1	0	2
97	0	1	2	0

Filter unnecessary data from beginning

WITH A AS

.

.

.

FROM A P1, A P2, A P3, MAX_TWO

WHERE P1.NA + P2.NA + P3.N = 100

AND P3.N + P2.NA <= 100

.

.

.

WHERE P1.NA + P2.NA + P3.NA + P4.N = 100

AND P4.N + P3.NA <= 100

AND P4.N + P3.NA + P2.NA <= 100

.

.

.

WHERE P1.NA + P2.NA + P3.NA + P4.NA + P5.N = 100

AND P5.N + P4.NA <= 100

AND P5.N + P4.NA + P3.NA <= 100

AND P5.N + P4.NA + P3.NA + P2.NA <= 100

.

.

.

Filter unnecessary data from beginning

- The principle of the minimum first step resultset
- If the conditions that can be put advance then do not put it behind
- Make full use of the ACCESS ability of Indexes
- Do not ignore the FILTER ability of Indexes

Summary

- Dealing with the problem as a whole, avoiding single row operation
- The principle of the minimum first step result set, reasonable choice of driving table
- Use new features like analysis functions to avoid repeated scans and reduce self-join
- Filter out unwanted data as much as possible at each step
- Write more SQL: Practice makes perfect
- Thinking more: The algorithm is key
- Perseverance: Optimization is endless

1 Introduction

Global supplier of end-to-end solutions of data assets

2 Solutions

Cloud Databases, Tuning, Operation and Maintenance, Recovery, consulting (Oracle & Mysql)

3 Products

zData/ MZ3/ZONE/
Yun mirror/DBPaaS

4 Customer&Cases

5000+ database
platforms
500+ Enterprises and
users in the industry

Well established in Chinese market, Enmotech keeps providing customers with all-around enterprise-class solutions that center data and talent with our globally top-ranking technical personnel (the directors of Oracle ACE, Oracle ACE experts, Oracle OCM experts, DB2 experts, middleware experts, and Unix system experts) and unique service ideas.



Eygle 盖国强

Oracle ACE Director

Guoqiang, Gai
Founder of ACOUG and
Enmotech.
The First Oracle ACE &
ACED in China



Kamus 张乐奕

Oracle ACE Director

Co-founder of ACOUG and
Enmotech, Extensive
hands-on experience with
Oracle Database RAC and
HA solutions, expertise in
fault diagnosis and
performance tuning



YangTingkun 杨廷琨

Oracle ACE Director

Enmotech CTO , expertise
in fault diagnosis and
performance tuning, and
known as Oracle
Encyclopedia



Xiong Jun 熊军

Oracle ACE Director

The core expert of ACOUG,
be good at trouble
shooting and data
recovery. The author of
the data recover software
Oracle ODU



Secooler 侯圣文

Oracle ACE Director

President of Enmoedu, he has Rich experience in Oracle, MySQL, Hadoop training courses, good at combining theory with practice teaching, easy to understand. As the gold medal lecturer of OCM, he has trained more than 600 OCMs and 10,000 DBAs.



Joel Perez

Oracle ACE Director

Joel is an Expert DBA (Oracle ACE Director, OCM Cloud Admin. & OCM11g) with over 15 years of Real World Experience in Oracle Technology, specialized in design and implement solutions of: High Availability, Disaster Recovery, Upgrades, Replication, Tuning, Cloud and all area related to Oracle Databases.



Roger 李真旭

Oracle ACE

Core expert of Enmotech, Good at Oracle database failure recovery, performance optimization and troubleshooting



Li yinan 李轶楠

Oracle ACE

Good at Oracle database application requirements analysis, architecture design, data modeling, database disaster management, performance optimization.

Customer & Cases & Services



- 31 Provinces
- 5 main filed (Financial Services、Communications operations、Energy Transportation、Government / public utility、Manufacturing / Commerce)
- 17+ Industry
- 500+ Enterprises and users in the industry
- 5000+ database platforms



Singapore

Hong Kong

Taiwan

Websites



<http://enmotech.com/>

<http://en.enmotech.com/>

<http://blog.enmotech.com/>



Thanks for listening



Contact us

Business Cooperation
marketing@enmotech.com
Business Consulting : (+86)
01059007017-7030



云和恩墨
ENMOTech

数据驱动 成就未来
Make Your Data Dance